

# ECE4703: Lecture 8

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# FFT Complexity Analysis: $N = 1$

DFT equation:

$$X[k] = \sum_{n=0}^{N-1} x[n]e^{-j2\pi kn/N} \quad \text{for } k = 0, 1, \dots, N - 1$$

When  $N = 1$ , the FFT and the DFT do the same computations and we simply have

$$X[0] = x[0]e^0$$

No multiplies or additions. Hence a one-point DFT/FFT has no MACs.

# FFT Complexity Analysis: $N = 2$

Two-point FFT computed using radix-2 decimation in time:

$$X[0] = X_{even}[0] + e^{-j2\pi \cdot 0/2} X_{odd}[0] = X_{even}[0] + X_{odd}[0]$$

$$X[1] = X_{even}[1] + e^{-j2\pi \cdot 1/2} X_{odd}[1] = X_{even}[1] - X_{odd}[1]$$

Note that  $X_{even}[k]$  and  $X_{odd}[k]$  for  $k = 0, 1$  are just one-point FFTs.

$$X_{even}[0] = \text{DFT}\{x[0]\} = x[0]$$

$$X_{odd}[0] = \text{DFT}\{x[1]\} = x[1]$$

$$X_{even}[1] = \text{DFT}\{x[0]\} = x[0]$$

$$X_{odd}[1] = \text{DFT}\{x[1]\} = x[1]$$

Hence,

$$X[0] = x[0] + x[1]$$

$$X[1] = x[0] - x[1]$$

How many MACs in the R2-DIT two-point FFT?

# FFT Complexity Analysis: $N = 4$ (part 1)

Four-point FFT computed using radix-2 decimation in time:

$$X[0] = X_{even}[0] + e^{-j2\pi \cdot 0/4} X_{odd}[0] = X_{even}[0] + X_{odd}[0]$$

$$X[1] = X_{even}[1] + e^{-j2\pi \cdot 1/4} X_{odd}[1] = X_{even}[1] - jX_{odd}[1]$$

$$X[2] = X_{even}[2] + e^{-j2\pi \cdot 2/4} X_{odd}[2] = X_{even}[2] - X_{odd}[2]$$

$$X[3] = X_{even}[3] + e^{-j2\pi \cdot 3/4} X_{odd}[3] = X_{even}[3] + jX_{odd}[3]$$

Remarks:

- ▶  $X_{even}[k]$  and  $X_{odd}[k]$  for  $k = 0, 1, 3, 4$  are all two-point FFTs.
- ▶ Recall that an  $N$ -point DFT/FFT is periodic in the sense that  $X[k] = X[k + N]$ .
- ▶ These facts imply that  $X_{even}[2] = X_{even}[0]$ ,  $X_{even}[3] = X_{even}[1]$ ,  $X_{odd}[2] = X_{odd}[0]$ , and  $X_{odd}[3] = X_{odd}[1]$ .
- ▶ Hence, we only need to compute four two-point FFTs:  $X_{even}[0]$ ,  $X_{even}[1]$ ,  $X_{odd}[0]$ , and  $X_{odd}[1]$ .

# FFT Complexity Analysis: $N = 4$ (part 2)

Four-point FFT computed using radix-2 decimation in time:

$$X[0] = X_{even}[0] + e^{-j2\pi \cdot 0/4} X_{odd}[0] = X_{even}[0] + X_{odd}[0]$$

$$X[1] = X_{even}[1] + e^{-j2\pi \cdot 1/4} X_{odd}[1] = X_{even}[1] - jX_{odd}[1]$$

$$X[2] = X_{even}[2] + e^{-j2\pi \cdot 2/4} X_{odd}[2] = X_{even}[2] - X_{odd}[2]$$

$$X[3] = X_{even}[3] + e^{-j2\pi \cdot 3/4} X_{odd}[3] = X_{even}[3] + jX_{odd}[3]$$

Computation of  $X_{even}[0]$  and  $X_{even}[1]$  requires how many MACs? 2

Computation of  $X_{odd}[0]$  and  $X_{odd}[1]$  requires how many MACs? 2

How many more MACs are required to assemble the two two-point FFTs into a four-point FFT? 4

Hence, the total MACs needed to compute a four-point FFT is 8.

# FFT Complexity Analysis: $N = 8$

Eight-point FFT computed using radix-2 decimation in time follows the same accounting:

$$X[k] = X_{even}[k] + e^{-j2\pi \cdot k/8} X_{odd}[k] \quad \text{for } k = 0, 1, \dots, 8$$

Computation of  $X_{even}[k]$  for  $k = 0, 1, 2, 3$  requires how many MACs? 8

Computation of  $X_{odd}[k]$  for  $k = 0, 1, 2, 3$  requires how many MACs? 8

How many more MACs are required to assemble the two four-point FFTs into an eight-point FFT? 8

Hence, the total MACs needed to compute an eight-point FFT is 24.

# FFT Complexity Analysis: General $N$

OPERATION	MACs
Computation of $N$ one-point FFTs	0
Assembling $\frac{N}{2}$ two-point FFTs from $N$ one-point FFTs	$N$
Assembling $\frac{N}{4}$ four-point FFTs from $\frac{N}{2}$ two-point FFTs	$N$
Assembling $\frac{N}{8}$ eight-point FFTs from $\frac{N}{4}$ four-point FFTs	$N$
$\vdots$	$\vdots$
Assembling one $N$ -point FFT from two $\frac{N}{2}$ -point FFTs	$N$

Hence, the total MACs needed to compute an  $N$ -point FFT is \_\_\_\_\_.