

# ECE504: Lecture 9

D. Richard Brown III

Worcester Polytechnic Institute

10-Nov-2009

## Lecture 10 Major Topics

We are finishing up Part II of ECE504: **Quantitative and qualitative analysis of systems**

mathematical description → results about behavior of system

Today, we will cover:

- ▶ Reachability of CT systems
- ▶ Controllability of CT systems
- ▶ Observability of CT systems
- ▶ Introduction to realization theory
- ▶ Minimal realizations

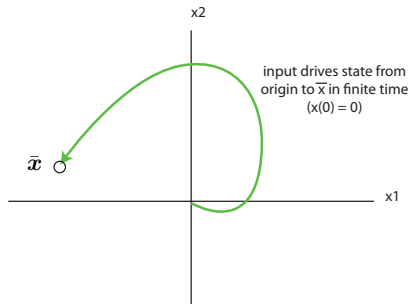
You should be reading Chen Chapters 6 and 7 now. Sections 6.1-6.3 discuss controllability and observability for CT systems. Sections 6.6-6.7 discuss controllability and observability for DT systems. Sections 7.1-7.2 give an introduction to realization theory.

# Reachable States (CT Systems)

## Definition

The state  $\bar{x} \in \mathbb{R}^n$  is a **reachable state** if there exists  $0 < T < \infty$  and an input function  $u(t)$  defined on  $t \in [0, T]$  such that  $x(T) = \bar{x}$  when  $x(0) = \mathbf{0}$  and when you apply the chosen input function  $u(t)$ .

The idea of reachability in CT systems is the same as in DT systems. We want to drive the state from  $\mathbf{0}$  to  $\bar{x}$  in finite time.



# Reachability Matrix/Theorem for CT-LTI Systems

## Definition

Given a CT-LTI system described by the matrices  $A$ ,  $B$ ,  $C$ , and  $D$ , the **reachability matrix** of this system is the matrix

$$\mathbf{Q}_r = [B \ AB \ \dots \ A^{n-1}B]$$

The reachability matrix for CT-LTI systems is defined in exactly the same way as for DT-LTI systems. The reachability theorem is also identical:

## Theorem

$\bar{x}$  is a reachable state if and only if  $\bar{x} \in \text{range}(\mathbf{Q}_r)$ .

The proof for this result is not as easy as the DT-LTI case, but is illustrative nonetheless...

# Proof Strategy — Preliminaries

First note that

$$[\exp\{t\mathbf{A}\}]^\top = \left[ \sum_{k=0}^{\infty} \frac{t^k}{k!} \mathbf{A}^k \right]^\top = \sum_{k=0}^{\infty} \frac{t^k}{k!} [\mathbf{A}^\top]^k = \exp\{t\mathbf{A}^\top\}$$

Now define the “ $T$ -reachability Grammian” corresponding to  $\mathbf{A}$  and  $\mathbf{B}$

$$\mathbf{M}_r(T) := \int_0^T \exp\{(T - \tau)\mathbf{A}\} \mathbf{B} \mathbf{B}^\top \exp\{(T - \tau)\mathbf{A}^\top\} d\tau$$

What are the dimensions of  $\mathbf{M}_r(T)$ ? \_\_\_\_\_

Does  $\mathbf{M}_r(T)$  have any special properties?

# Proof Strategy

We've defined

$$\mathbf{M}_r(T) := \int_0^T \exp\{(T - \tau)\mathbf{A}\} \mathbf{B} \mathbf{B}^\top \exp\{(T - \tau)\mathbf{A}^\top\} d\tau$$

$$\mathbf{Q}_r = [\mathbf{B} \ \mathbf{A}\mathbf{B} \ \dots \ \mathbf{A}^{n-1}\mathbf{B}]$$

The strategy: For arbitrary  $0 < T < \infty$ , we will show

1.  $\text{range}(\mathbf{M}_r(T)) \stackrel{(a)}{\subset} \{\text{reachable states}\} \stackrel{(b)}{\subset} \text{range}(\mathbf{Q}_r)$
2.  $\text{range}(\mathbf{M}_r(T)) \stackrel{(c)}{=} \text{range}(\mathbf{Q}_r)$

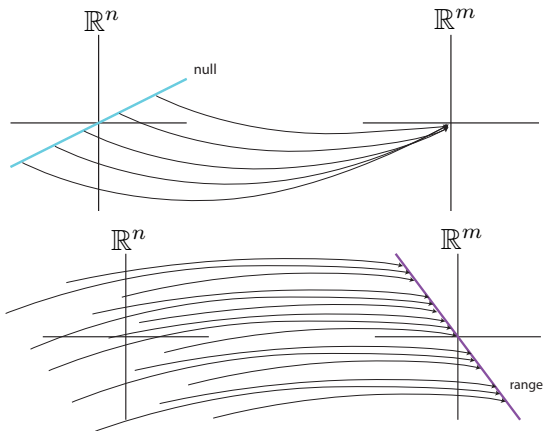
The second part implies the equality of two subspaces. By “sandwiching”, the third subspace (the set of reachable states) must also be equal to the other two subspaces.

Bonus: We can use the Grammian to test for reachability.

# Strategy Part 1 Proof (show (a) and (b))

# Strategy Part 2 Proof — Preliminaries

This part is going to require two new linear algebra results. First, recall that, for  $\mathbf{A} \in \mathbb{R}^{n \times m}$ ,  $\text{range}(\mathbf{A})$  is a subspace of  $\mathbb{R}^n$  and  $\text{null}(\mathbf{A})$  is a subspace of  $\mathbb{R}^m$ . What can you say about  $\text{range}(\mathbf{A}^\top)$  and  $\text{null}(\mathbf{A}^\top)$ ?





## Strategy Part 2 Proof — Preliminaries

Here are our two new linear algebra results:

### Lemma

$\dim(\text{range}(\mathbf{A})) + \dim(\text{null}(\mathbf{A}^\top)) = n$  for any  $\mathbf{A} \in \mathbb{R}^{n \times m}$ .

### Lemma

If  $\mathbf{v} \in \text{range}(\mathbf{A})$  and  $\mathbf{w} \in \text{null}(\mathbf{A}^\top)$ , then  $\mathbf{v}^\top \mathbf{w} = 0$ .

What does this imply about  $\text{range}(\mathbf{A})$  and  $\text{null}(\mathbf{A}^\top)$ ?

# Strategy Part 2 Proof (show (c))

# Remarks on the Reachability Theorem

We now know that the set of reachable states is equal to  $\text{range}(\mathbf{Q}_r)$  and is a subspace of  $\mathbb{R}^n$ .

**Interesting consequence:** The reachability definition requires that we drive the state from the origin to  $\bar{x}$  in a finite time. The reachability theorem implies that we can always drive the state from  $\mathbf{0}$  to  $\bar{x} \in \text{range}(\mathbf{Q}_r)$  **as quickly as we want** in CT-LTI systems.

## Theorem

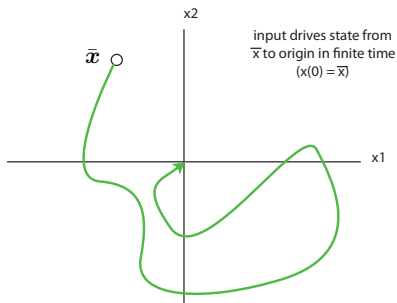
*The set of reachable states  $\text{range}(\mathbf{Q}_r)$  is invariant under  $\mathbf{A}$ , i.e. if  $\mathbf{x} \in \text{range}(\mathbf{Q}_r)$  then  $\mathbf{Ax} \in \text{range}(\mathbf{Q}_r)$ .*

The proof is a consequence of the Cayley-Hamilton theorem.

# Controllability versus Reachability for CT Systems

## Definition

The state  $\bar{x} \in \mathbb{R}^n$  is a **controllable state** if there exists  $0 < T < \infty$  and an input function  $u(t)$  defined on  $t \in [0, T]$  such that  $x(T) = \mathbf{0}$  when  $x(0) = \bar{x}$  and when you apply the chosen input function  $u(t)$ .



What can we say about the relationship between reachable states and controllable states for CT systems?

# More Remarks on the Reachability Theorem

## Lemma

*$\bar{x}$  is a controllable state if and only if  $\bar{x}$  is a reachable state.*

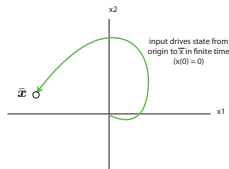
The set of reachable states is identical to the set of controllable states in CT-LTI systems **because the CT-STM is always invertible**. This is not true, as we saw, for DT-LTI systems.

## Definition

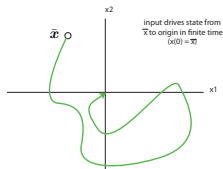
A CT-LTI system with  $\text{range}(\mathbf{Q}_r) = \mathbb{R}^n$ , i.e. all states are reachable (or controllable), is called a “reachable” (or “controllable”) system.

# Summary: Reachability/Controllability for CT-LTI Systems

- ▶ Reachability: Drive state from  $\mathbf{0}$  to  $\bar{\mathbf{x}}$ .



- ▶ Controllability: Drive state from  $\bar{\mathbf{x}}$  to  $\mathbf{0}$ .



- ▶  $\{\text{reachable states}\} = \text{range}([\mathbf{B} \ \mathbf{A}\mathbf{B} \ \dots \ \mathbf{A}^{n-1}\mathbf{B}]) = \text{range}(\mathbf{Q}_r)$ .
- ▶  $\{\text{reachable states}\} = \{\text{controllable states}\}$ .
- ▶ If  $\bar{\mathbf{x}}$  and  $\tilde{\mathbf{x}}$  are both reachable/controllable then we can find an input to drive the system from  $\bar{\mathbf{x}}$  to  $\tilde{\mathbf{x}}$  (or vice-versa) as quickly as we want.

# Observability for CT-LTI Systems

## Definition

The state  $\bar{x} \in \mathbb{R}^n$  is an **unobservable state** if, for any choice of input function  $\mathbf{u}(t) : [0, \infty) \mapsto \mathbb{R}^p$ , the output function  $\mathbf{y}(t) : [0, \infty) \mapsto \mathbb{R}^q$  given initial state  $\mathbf{x}(0) = \bar{x}$  is the same as the output function given initial state  $\mathbf{x}(0) = \mathbf{0}$ .

Intuition:

- ▶ Basically the same idea as DT-LTI systems.
- ▶ You want to determine the initial state of the system  $\mathbf{x}(0)$ .
- ▶ You are allowed to apply any input you want to the system and measure the resulting output.
- ▶ The state  $\mathbf{x}(0) = \bar{x}$  is called “unobservable” if you are unable to distinguish it from the initial state  $\mathbf{x}(0) = \mathbf{0}$  for any input.

# Observability for CT-LTI Systems

Given  $\mathbf{x}(0) = \bar{\mathbf{x}}$ , the output of a CT-LTI system is

$$\mathbf{y}(t) = \mathbf{C} \exp\{t\mathbf{A}\}\bar{\mathbf{x}} + \int_0^t \mathbf{C} \exp\{(t - \tau)\mathbf{A}\}\mathbf{B}\mathbf{u}(\tau) d\tau + \mathbf{D}\mathbf{u}(t)$$

Given  $\mathbf{x}(0) = \mathbf{0}$ , the output of a CT-LTI system is

$$\mathbf{y}(t) = \int_0^t \mathbf{C} \exp\{(t - \tau)\mathbf{A}\}\mathbf{B}\mathbf{u}(\tau) d\tau + \mathbf{D}\mathbf{u}(t)$$

Hence, for these to be equal, we must have

$$\boxed{\mathbf{C} \exp\{t\mathbf{A}\}\bar{\mathbf{x}} = \mathbf{0} \quad \forall t \geq 0}$$

This condition is equivalent to the statement that “ $\bar{\mathbf{x}}$  is unobservable”.



# Observability Matrix/Theorem for CT-LTI Systems

## Definition

Given a CT-LTI SS system described by the matrices  $A$ ,  $B$ ,  $C$ , and  $D$ , the **observability matrix** of the system is defined as

$$Q_o = \begin{bmatrix} C \\ CA \\ \vdots \\ CA^{n-1} \end{bmatrix}$$

Note that this is exactly the same definition as we saw for DT-LTI systems.

## Theorem

$\bar{x}$  is an unobservable state if and only if  $\bar{x} \in \text{null}(Q_o)$ .

Proof...

# Observable Systems

## Definition

A CT-LTI system is **observable** if no  $\bar{x} \neq \mathbf{0}$  is an unobservable state.

In other words, like DT-LTI systems, a CT-LTI system is observable if  $\dim(\text{null}(\mathbf{Q}_o)) = 0$  or, equivalently, if  $\text{rank}(\mathbf{Q}_o) = n$ .

Some possibly useful facts:

- ▶ The set of unobservable states is a subspace of  $\mathbb{R}^n$ .
- ▶ This subspace is invariant under  $\mathbf{A}$ , i.e. if  $\bar{x}$  is in the set of unobservable states, then so is  $\mathbf{A}\bar{x}$  (consequence of the Cayley-Hamilton theorem).

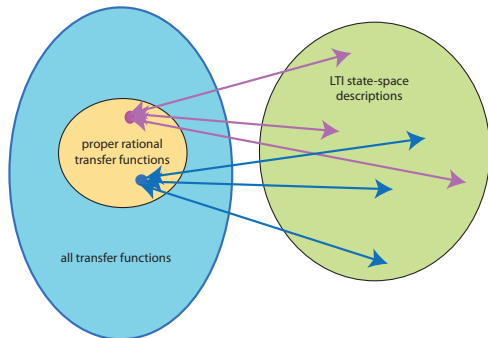
# Final Remarks on Reachability, Controllability, and Observability

1. Reachability and Controllability are “control” concepts: “How can we affect the behavior of a system through the input”? This is a preview of Part III of the course.
2. Chapter 6 of Chen has many more results (and examples) than covered in the lecture.
3. Chen does not distinguish between reachability and controllability for CT systems.
  - 3.1 Remember: The set of reachable states and the set of controllable states are the same for CT-LTI systems but **not necessarily the same for DT-LTI systems**.
  - 3.2 The first mention of reachability is in Section 6.6.1.
4. Please see the discussion of reachability, controllability, and observability for LTV systems in Section 6.8.

# Introduction to Realization Theory

Review questions:

1. You are given a  $p$ -input,  $q$ -output transfer function  $\hat{G}(s)$ . Can you always find a state-space description for this system?
2. You are given a  $p$ -input,  $q$ -output,  $n$ -state LTI state-space description  $\{A, B, C, D\}$ . Can you always find a transfer function?



# Introduction to Realization Theory

Recall that

$$\hat{G}(s) = C(sI_n - A)^{-1}B + D$$

- ▶ All LTI state-space systems can be converted to transfer functions.
- ▶ All transfer functions are not “realizable”.
- ▶ Any realizable transfer function has infinitely many realizations. For example, if  $\{A, B, C, D\}$  is a realization, then for any invertible  $P$ ,  $\{P^{-1}AP, P^{-1}B, CP, D\}$  is also a realization.
- ▶ If a system with  $n$  states realizes the transfer function  $\hat{G}(s)$ , then we can always create a system with more than  $n$  states that also realizes  $\hat{G}(s)$ . Example...

# Minimal Realizations

Suppose  $\hat{G}(s)$  is realizable. How small can the state dimension  $n$  be?

## Definition

Given an CT-LTI system with realizable transfer function  $\hat{G}(s)$ , the **McMillan degree** of the system is the smallest possible state vector dimension over the class of all  $\{A, B, C, D\}$  that realize the system with transfer function  $\hat{G}(s)$ .

## Definition

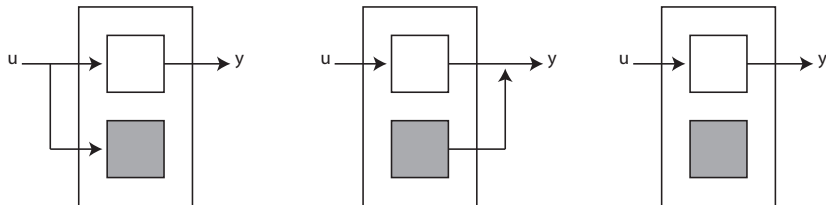
A realization  $\{A, B, C, D\}$  is **minimal** if its state vector dimension equals the McMillan degree.

Question: Is the minimal realization unique?

# Non-Minimal Realizations: Intuition

The system isn't minimal if we have extraneous states.

There are three ways to have extraneous states:



In each case, the states in the shaded box have no impact on the I/O behavior of the system.

# Fundamental Theorem of Realization Theory

## Theorem

A realization  $\{A, B, C, D\}$  is **minimal** if and only if  $A$  and  $B$  are such that the system is reachable and  $A$  and  $C$  are such that the system is observable.

This should be intuitively satisfying.

- ▶ If a system is reachable, then all states are connected to the input.
- ▶ If a system is observable, then all states are connected to the output.

For a system to be minimal, **all states must be connected to both the input and the output.**

Proof...



# Conclusions

Today we covered:

- ▶ Reachability for CT systems.
- ▶ Controllability for CT systems.
- ▶ Observability for CT systems.
- ▶ Introduction to realization theory
- ▶ Minimal realizations

Next time:

- ▶ More realization theory.
- ▶ Start Part III of course: Controlling the behavior of systems.